

# Ch. 17 Problems

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14-6

- (a) Long shaping times  $\Rightarrow$  accurate measurement of collected charge  
 $\rightarrow$  Long shaping times give better pulse height resolution.
- (b) From fig. 17-14, we can see that the best signal: noise ratios occur for monopolar shaping methods. Since random noise degrades amplitude (pulse height) information, we want the best signal: noise ratio possible. Thus, choose monopolar shaping over bipolar for better pulse height resolution.
- (c) Baseline restoration introduces noise to the system, thus decreasing the signal: noise ratio. Although its benefits may outweigh its problems for high rate systems, here we are dealing with a low average rate (as stated in the problem), so if we want the best preservation of pulse height information we should not use active baseline restoration.

19-9

$$f_e = \frac{e^{-nZ}}{nZ + 1}$$

$$n = \text{total count rate} = 25,000 \text{ s}^{-1}$$

$$Z = 4 \times 10^{-6} \text{ s}$$

$$\therefore f_e = \frac{\exp[0.1]}{0.1 + 1} \quad (\text{non-paralyzable})$$

$$f_{e, NP} = 82.3\% \quad (\text{fraction of events escaping pileup})$$

$$f_e = e^{-2nZ} \quad (\text{paralyzable})$$

$$= e^{-2(1.1)}$$

$$f_{e, P} = 81.9\%$$

~~fraction of total events lost is the same as the fraction of full energy events lost.~~

Since pileup can occur with any pulse, the fraction of total events lost is the same as the fraction of full energy events lost.

$$\text{Also, } f_{\text{lost}} = 1 - f_e$$

$$\Rightarrow f_{\text{lost}}^{(\text{paralyzable})} = 18.1\%$$

$$f_{\text{lost}}^{(\text{non-paralyzable})} = 17.7\%$$

17-11 Chance coincidence rate =  $2\tau\Gamma_1\Gamma_2$

Since the pulses are random + uncorrelated, all coincidences will be chance coincidences.

Anti-coincidence units provide an output pulse each time there is a coincidence event (so that the unit can only supply an output when there is not an anti-coincidence pulse). Thus the anti-coincidence output rate is just the coincidence rate which here is given by the chance coincidence rate:

$\Gamma_{ch} = 2\tau\Gamma_1\Gamma_2$  ✓

17-13  $t_{delay} = l/v_{cc}$

where  $l$  = cable length  
 $v_{cc}$  = cable carrier velocity  $\approx c = 3 \times 10^8 \text{ m/s} = 0.3 \text{ m/ns}$

$\Rightarrow l = t_{delay} \cdot 0.3 \text{ m/ns}$    
not true in a cable use Table in text

$l = (100 \text{ ns}) (0.3 \text{ m/ns}) * \text{velocity fraction} \sim 0.659$

$l = 30 \text{ m}$

→ corrected on next page.

$$\boxed{7-13} \quad t_{\text{delay}} = l/v_{cc}$$

$l$  = cable length

$$v_{cc} = \text{current carrier velocity} = (3 \times 10^8 \text{ m/s}) [0.659] = \underline{\underline{0.1977 \text{ m/ns}}}$$

from table 16-1  
for a polyethylene  
coaxial cable.

$$\begin{aligned} \Rightarrow l &= t_{\text{delay}} v_{cc} \\ &= (100 \text{ ns}) (0.1977 \text{ m/ns}) \end{aligned}$$

$$\therefore \boxed{l = 19.8 \text{ m}}$$

17-15

$$I_c = \epsilon_1 \epsilon_2 S$$

$$R_{ch} = 2\pi r_1 r_2$$

$$r_1 = \epsilon_1 S$$

$$r_2 = \epsilon_2 S$$

$$\Rightarrow r_1 r_2 = \epsilon_1 \epsilon_2 S^2$$

$$\downarrow$$
$$R_{ch} = 2\pi \epsilon_1 \epsilon_2 S^2 \Rightarrow$$
$$R_t = \epsilon_1 \epsilon_2 S$$

$$\frac{R_t}{R_{ch}} = \frac{1}{2\pi S}$$

Ⓐ Changing solid angle changes  $\epsilon_1$  and  $\epsilon_2$   
 $\Rightarrow$  changing solid angle has no effect on  $R_t/R_{ch}$

$\rightarrow$  Practical considerations: limited by geometry of the detectors

Ⓑ Varying source material varies  $S$

$\Rightarrow$  less source material = increased  $R_t/R_{ch}$

$\rightarrow$  still need enough source material to get good statistics

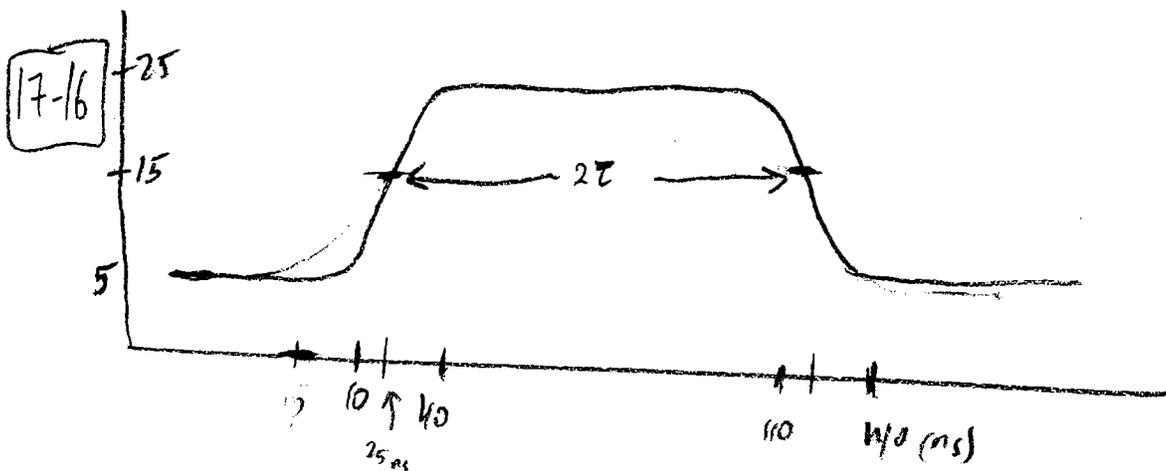
Ⓒ  $R_t/R_{ch} \propto 1/\tau \Rightarrow$  lowering  $\tau$  increases  $R_t/R_{ch}$ .

$\rightarrow$  need  $\tau$  to be at least as large as  $1/2$  the width (at the base) of the prompt coincidence peak, otherwise we cut out true coincidences.

Ⓓ Varying the pulse height window just varies the rates  $r_1$  and  $r_2$ : it is the same as having the efficiencies  $\epsilon_1$  and  $\epsilon_2$  change. Since  $R_t/R_{ch}$  has no dependence on

$\epsilon_1$  and  $\epsilon_2$ , varying pulse height window has no effect on  $R_t/R_{ch}$ .

$\rightarrow$  can't make the window too restrictive or else we will cut out too many events



(a)  $2\tau$  = FWHM of curve from 10 ns  $\rightarrow$  110 ns.

curve is approx. linear from 10 ns  $\rightarrow$  40 ns and 110 ns  $\rightarrow$  140 ns

(0-20): Slope =  $\frac{20}{30} = \frac{2}{3}$

$\Rightarrow f(t) = 15 = \frac{2}{3}t - \frac{5}{3}$

offset =  $-\frac{5}{3}$

$\Rightarrow t_1 = 25$  ns

$t_2 = 125$  ns

FWHM =  $t_2 - t_1 = 125 - 25 = 100$  ns

$\Rightarrow 2\tau = 100$  ns  $\Rightarrow \tau = 50$  ns ✓

Corrections on next page

(b) Width of <sup>Print</sup> coincidence peak =  $2\tau = 100$  ns

no it shows up in rise time

(c)  $t_{c1} = 2\tau, t_{c2} = 5$  (ns) <sup>per second on graph</sup> (from graph)

$\nu_1 = \nu_2 \Rightarrow \nu = \left(\frac{5}{2\tau}\right)^{1/2} = \left(\frac{5 \text{ ns}}{2(50 \text{ ns})}\right)^{1/2} = \sqrt{0.05} \text{ ns}^{-1} \Rightarrow$

$\nu = 0.224 \text{ ns}^{-1}$

(K-16) (b) here  $\tau$  is longer than  $\frac{1}{2}$ -the width (at the base) of the prompt coincidence peak (because of the flat-topped plateau)

$\Rightarrow$  width of prompt coincidence peak will show up within the plateau,

- Use the marks on the graph as the width;

$$\Rightarrow W = 110_{\text{ns}} - 40_{\text{ns}}$$

$$W = 70_{\text{ns}}$$

$$(c) \Gamma_{\text{ch}} = 2\tau\Gamma_1\Gamma_2 = 5 \text{ s}^{-1}$$

$$\Gamma_1 = \Gamma_2 \equiv \Gamma \Rightarrow \Gamma = \left[ \frac{5 \times 10^{-9} (\text{ns}^{-1})}{2\tau} \right]^{1/2} = \left[ \frac{5 \text{ s}^{-1}}{2(50 \times 10^{-9} \text{ s})} \right]^{1/2} = \sqrt{5 \times 10^7 \text{ s}^{-1}}$$

$$\Rightarrow \Gamma = 7071 \text{ s}^{-1}$$

(17-17) # of comparators for  $n$  bits =  $2^n - 1$

$$\Rightarrow \left\{ \begin{array}{l} 3 \text{ stages of 4 bits each} \Rightarrow (2^4 - 1) \cdot 3 = \underline{45 \text{ comparators}} \\ 4 \text{ " " " 3 " " " } \Rightarrow (2^3 - 1) \cdot 4 = \underline{28 \text{ comparators}} \end{array} \right.$$

$\Rightarrow$  4 stages of 3 bits each requires less total comparators.

Latency Time  $\propto$  # of <sup>stages</sup> ~~comparators~~ =  $2^{\text{\# of stages}}$

$\Rightarrow$  3 stages of 4 bits each has lower latency time.

$$\left( \underline{2^3 = 8} \text{ - VS - } \underline{2^4 = 16} \right)$$